

The *G*rid Method in Numerical \mathbb{R} Real Algebraic Geometry

A short presentation at the XXI Santaló School of Mathematics

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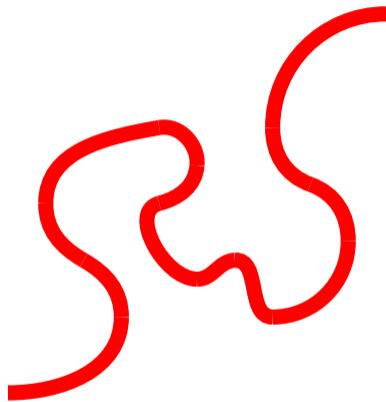
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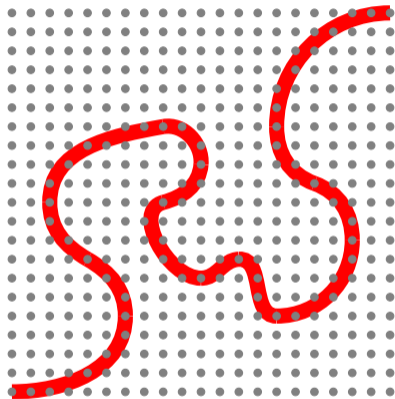
What is the *G*rid Method?

What is the \mathcal{G} rid Method? I



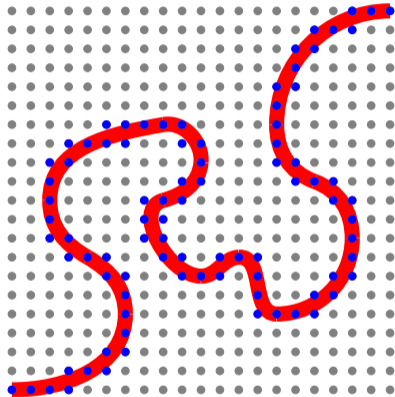
We want to capture a geometric object

What is the \mathcal{G} rid Method? I



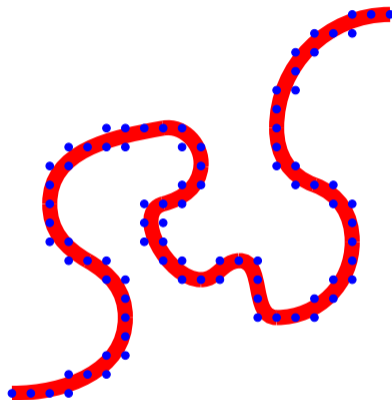
We cover the ambient space by a grid

What is the \mathcal{G} rid Method? I



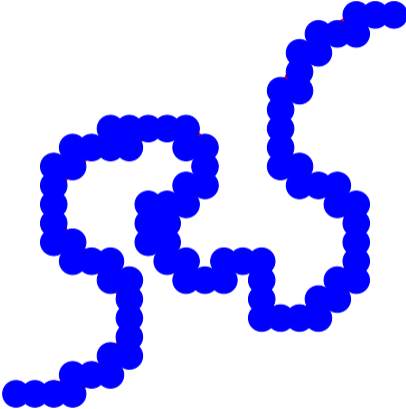
We select the nearby point to approximate the object

What is the \mathcal{G} rid Method? I



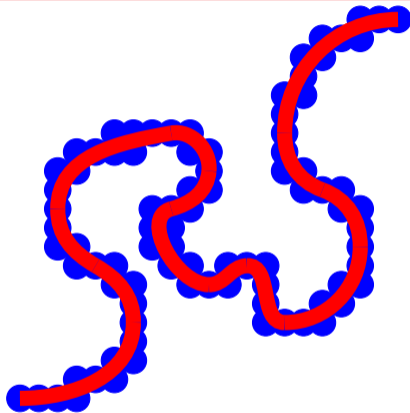
We select the nearby point to approximate the object
...and forget the rest!

What is the Grid Method? I



We postprocess the selection

What is the \mathcal{G} rid Method? I



We postprocess the selection
to capture what we want about the geometric object of interest

What is the *G*rid Method? II

1. Cover with a grid the ambient space where the desired object lies
 - Which properties should the covering grid have?
 - How can we generate such a covering grid efficiently?
2. Select the points that approximate the desired object
 - What means to 'approximate' the desired object?
 - How do we select the points so that we approximate the object?
 - How to refine the grid when selection fails? (Non-Adaptive vs. Adaptive)
3. Postprocess the points to obtain the information about the desired object
 - How can I combine the information coming from the selected points?
 - Can I do so in a fast way?

Numerical \mathbb{R} Real Algebraic Geometry (NRAG)

NRAG I: What is NRAG?

What do we deal with?

Problems defined by real polynomials
considering errors in the given input

E.g. feasibility—*is there a real zero?*—, counting real zeros,
Betti numbers/homology groups of a semialgebraic set. . .

NRAG II: The basics of condition numbers

Condition number:

- Measure of the sensitivity of the output to errors of the input for a specific problem
- Complexity of numerical algorithms depends on input size and condition number
 - same input-size doesn't mean similar run-times!
 - inputs with infinite condition number (*ill-posed*) cannot be handled numerically*
- Probabilistic analysis of the condition number for a random input
 - gives probabilistic complexity

*If input is assumed to be real. If input is assume to be integer, then. . .

NRAG III: Condition-based and prob. complexity

Worst-case complexity:

$$\max\{\text{run-time}(\text{ALGORITHM}, i) \mid \text{size}(i) \leq s\}$$

Condition-based complexity:

$$\max\{\text{run-time}(\text{ALGORITHM}, i) \mid \text{size}(i) \leq s, \text{cond}(i) \leq c\}$$

—the condition number allows to explain better the behaviour of **ALGORITHM** at input i

Probabilistic complexity:

$$\mathbb{E}\{\text{run-time}(\text{ALGORITHM}, i) \mid \text{size}(i) \leq s\}$$

...or randomly perturbed arbitrary input (smoothed paradigm)

$$\max \left\{ \mathbb{E} \left(\text{run-time}(\text{ALGORITHM}, \tilde{i}) \mid \text{dist}(\tilde{i}, i) \leq \sigma \right) \mid \text{size}(i) \leq s \right\}$$

NRAG IV: Local condition number

Let $f \in \mathcal{H}_{n,d}[q] := \prod_{i=1}^q \mathbb{R}[X_0, \dots, X_n]_{d_i}$ and $x \in \mathbb{S}^n$, the *local condition number* of f at x is

$$\kappa_W(f, x) := \frac{\|f\|_W}{\sqrt{\|f(x)\|_2^2 + \sigma_q(\Delta^{-1/2} D_x f)^2}}$$

where

- $\|\cdot\|_W$ is the Weyl norm—the Weyl norm is not the only choice!—,
- σ_q the q th singular value,
- Δ the diagonal matrix with d_1, \dots, d_q in the diagonal, and
- $D_x f$ the tangent map $T_x \mathbb{S}^n \rightarrow \mathbb{R}^q$ of f at x .

Main observation: $\kappa_W(f, x) = \infty$ iff x is a singular zero of f

NRAG V: Global condition number

Let $f \in \mathcal{H}_{n,d}[q] := \prod_{i=1}^q \mathbb{R}[X_0, \dots, X_n]_{d_i}$,

(G) The *global condition number* of f is

$$\kappa_W(f) := \max_{x \in \mathbb{S}^n} \kappa_W(f, x).$$

- Controls complexity of non-adaptive grid methods
- **Probabilistic properties:** Small with high probability, but infinite expectation

(A) The *global-average condition number* is

$$\kappa_W^{\text{av}}(f) := \sqrt[n]{\mathbb{E}_{x \in \mathbb{S}^n} \kappa_W(f, x)^n}.$$

- Controls complexity of adaptive grid methods
- **Probabilistic properties:** Finite moments up to order strictly less than $n + 1$

A Brief History of the *G*rid Method in NRAG

A Brief History of the \mathcal{G} rid Method in N \mathbb{R} AG

- (Cucker & Smale, 1999) **Feasibility of semialgebraic sets**
—no probabilistic analysis
- (Cucker, Krick, Malajovich & Wschebor; 2008, 2009 & 2011) **Counting real zeros**
—with high probability under Gaussian assumptions, but no finite expectation
—under robust assumptions by (Ergür, Paouris & Rojas; 2019 & 2021)
- (Cucker, Krick & Shub, 2018) **Homology of zero sets**
—with high probability under under Gaussian assumptions, but no finite expectation
—also under robust assumptions (T-C; unpublished)
- (Bürgisser, Cucker & Lairez; 2018) **Homology of basic semialgebraic sets**
—with high probability, but no finite expectation
- (Bürgisser, Cucker & T-C; 2020 & 2022) **Homology of general semialgebraic sets**
—with high probability, but no finite expectation
- (T-C; MEGA 2021 & to appear in arXiv in 2022) **Counting & computing real zeros**
—adaptively with finite expectation under robust assumptions!!!

Challenges and Future of the *G*rid Method in N_RAG

Challenges and Future of the \mathcal{G} rid Method in NRAG

- Construction of nice grids efficiently
 - we want to construct fast grids with structure we can exploit to compute faster (this is why I am here!)
- Postprocessing the subselection of the grid
 - this is one of the bottlenecks for using adaptive grids for homology computation
- Can we make our grids depend on the input?
 - the grids we construct are not input-sensitive
- Exploit the polynomial nature of the input
 - until now we only use that polynomials are C^2 -functions
 - recent progress on this! (keep attention to arXiv this year)

Eskerrik asko por su atención!